Maximum number of independence atoms in a relational database

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This research problem deals with some combinatorial question of relational databases. For our purposes the database is simply a matrix, whose columns are the *attributes*, the rows are the individual records of data, frequently *tuples*. For example consider Table 1, where the attributes are named.

Course Name	Year	Lecturer	Credits	Semester
Mathematics	2019	Cornelius	5	1
Datamining	2018	Sarah	7	2
Theory of Algorrithms	2019	Sarah	7	2

Table 1: A sample database table.

Independence and conditional independence are fundamental concepts in areas as diverse as artificial intelligence, probability theory, social choice theory, and statistics. In this problem we investigate a special case of independence, namely independence atoms. Intuitively, a relation (table) r satisfies the independence atom $X \perp Y$ between two disjoint sets X and Y of attributes, if for all tuples $t_1, t_2 \in r$ there is some tuple $t \in r$ which matches the values of t_1 on all attributes in X and matches the values of t_2 on all attributes in Y. In other words, in relations that satisfy $X \perp Y$, the occurrence of X-values is independent of the occurrence of Y-values.

Example 1. Consider a simple database schema that stores information about the enrolment of students into a fixed course. In fact, the schema records for each enrolled student, the year in which they completed a prerequisite course. More formally, we have the schema (set of attributes) Enrol=S(tudent), P(rerequisite), Y(ear). Intuitively, every student must have completed every prerequisite in some year. For this reason, for any value in the Student column and every value in the Prerequisite column there is some year for when this student has completed that prerequisite. That is, the values in the Student column are independent of the values in the Prerequisite column. A snapshot relation r over Prerequisite column are independent of the values in the Prerequisite column.

Student	Prerequisite	Year
Turing	Math201	1932
Gödel	Math201	1925
Turing	Phys220	1932
Gödel	Phys220	1925

A collection Σ of independence atoms imply independence atom σ , denoted by $\Sigma \models \sigma$, if every table that satisfies Σ , also satisfies σ . This implication can be described equivalently by derivation rules, as follows. Let X, Y, Z be subsets of attributes.

- 1. $X \perp \emptyset$ holds always.
- 2. $X \perp Y \Rightarrow Y \perp X$ (symmetry).
- 3. $X \perp YZ \Rightarrow X \perp Y$ (decomposition).
- 4. $X \perp Y \wedge XY \perp Z \Rightarrow X \perp YZ$ (exchange).

So we have that $\Sigma \models \sigma$ iff σ can be derived starting from independence atoms in Σ using a finite number of applications of the rules above. XY means the union of attribute sets X and Y, similarly YZ is the union of Y and Z. (In general if two sets of attributes are written next to each other then it denotes the union of the two sets. This is traditional notation in the area.)

The main concept is non-redundance of set of independence atoms. Σ is non-redundant iff no $\sigma \in \Sigma$ can be derived from $\Sigma \setminus \{\sigma\}$, that is no independence atom in Σ can be derived from the other atoms in Σ .

Research problem We want to find f(n) the largest size of a non-redundant collection of independence atoms in a schema of n attributes.

What do we know? One must observe that none of the derivation rules introduces new attributes, so $\{X_i \perp Y_i : |X_i \cup Y_i| = \lceil \frac{n}{2} \rceil, i = 1, 2, \dots, \binom{n}{\lceil \frac{n}{2} \rceil}, X_i \cup Y_i \neq X_j \cup Y_j \text{ for } i \neq j\}$ with none of X_i and Y_i being nonempty, is non-redundant. Thus, $f(n) \geq \binom{n}{\lceil \frac{n}{2} \rceil}$.

There is a divide-and-conquer algorithm that decides for a given collection Σ of independence atoms and another independence atom σ whether Σ implies σ . It is implemented in python, if someone wants to try it, write me an email.

We have better constructions then the one indicated above. This shows that $f(5) \ge 12 > {5 \choose 2}$.

Qualifying questions

- 1. Show that f(3) = 6.
- 2. Show that $f(5) \geq 11$.
- 3. Show that $f(5) \geq 12$.

4. Show that for pairwise disjoint attribute sets U, V, W, W', W'', Y the following implication holds. $UV \perp WW' \wedge UWW'' \perp Y \Rightarrow U \perp YW$.

Questions, solutions should be directly sent to saliattila@gmail.com $\,$